

Chapter 3: Process management part II

- Semaphore Definition
- P and V operation
- Mutex using P and V operation
- Process Synchronization
- Synchronization using P and V operation
- Typical Problems

Semaphore

- It including two parts
A integer ,its initial value ≥ 0
and a queue pointer

P and V operation on Semaphore

- P operation:

- ① $sem-1$
- ② If $sem \geq 0$ after $(sem-1)$, then it return, the process continue executing
- ③ If $sem < 0$ after $(sem-1)$, then the process be blocked and put to the block queue.

- V operation

- ① $sem+1;$

- ② If $sem > 0$ after $(sem+1)$, the process continue executing

- ③ If $sem \leq 0$ after $(sem+1)$, then awake a process in the blocking queue

How to implement P and V operation

Mutex using P and V operation

■ P1

P(sem)

<S>

V(sem)

P2

P(sem)

<S>

V(sem)

Synchronization

- An synchronization example
- Private and public Semaphore

Synchronization using P and V

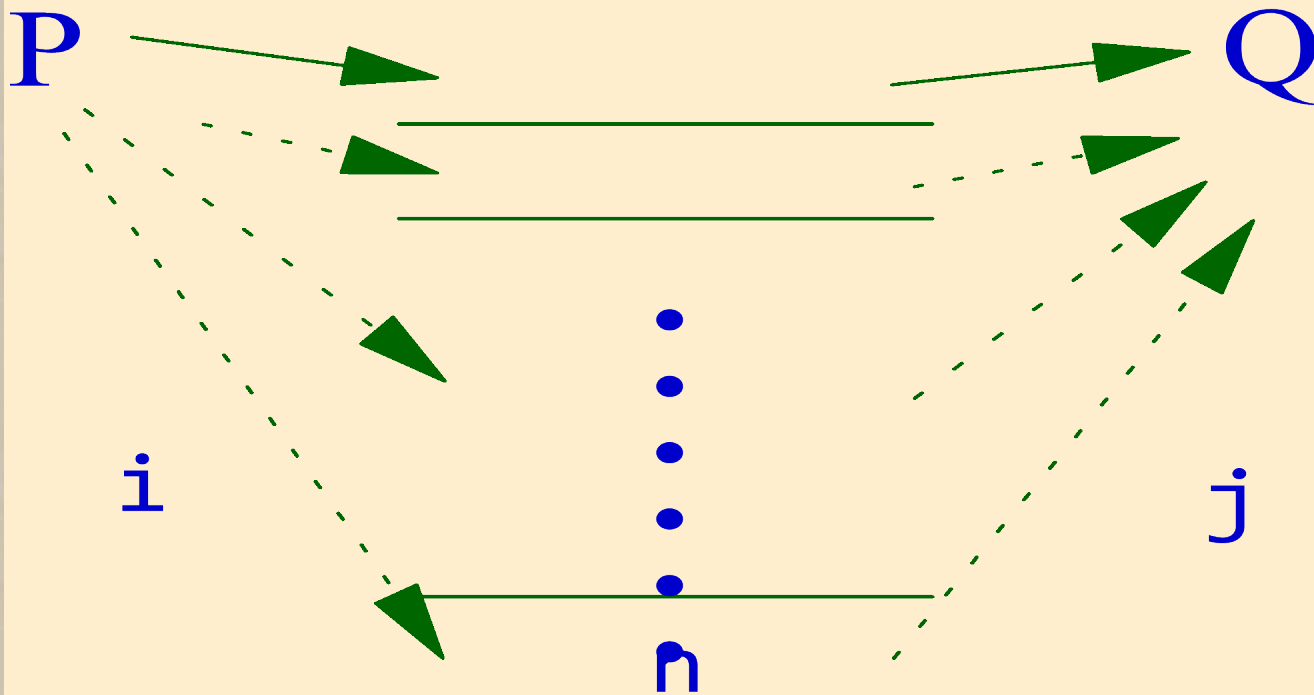
- Three steps:
 - ① Deciding how many semaphore should use
 - ② Initialize the semaphore
 - ③ Using P and V operation

Typical problems

- One producer and one consumer and one buffer
- Multi producers and multi consumers
- Barber shop problem
- Dining-Philosophers Problem

放消息

取消息



n 个缓冲区
(Buffer)

P:

```
i = 0;
while (true) {
    producing;
    P(S1);
    put into Buffer [i];
    V(S2);
    i = (i+1) % n;
};
```

Q:

```
j = 0;
while (true) {
    P(S2);
    take from Buffer[j];
    V(S1);
    consuming;
    j = (j+1) % n;
};
```

Multi producer and consumers

- Shared data

semaphore full, empty, mutex;

Initially:

full = 0, empty = n, mutex = 1

以上内容仅为本文档的试下载部分，为可阅读页数的一半内容。如要下载或阅读全文，请访问：<https://d.book118.com/146011113121010213>