

2022 新疆软件水平考试考试真题卷（4）

本卷共分为 1 大题 50 小题，作答时间为 180 分钟，总分 100 分，60 分及格。

一、单项选择题（共 50 题，每题 2 分。每题的备选项中，只有一个最符合题意）

1. 在 UML 提供的图中，（28）用于对系统的静态设计视图建模，（29）用于对系统的静态实现视图建模。（29）处填（）。

- A. 类图
- B. 对象图
- C. 部署图
- D. 构件图

2. 8 层完全二叉树至少有（52）个结点，拥有 100 个结点的完全二叉树的最大层数为（53）。（53）处填（）。

- A. 6
- B. 7
- C. 8

D. 9

3. 广义表 $(a, (a, b), d, e, ((i, j), k))$ 的长度是 (48) , 深度是 (49) 。(49) 处填 ()。

A. 3

B. 4

C. 5

D. 6

4. 寻址是指控制器根据指令的地址码寻找操作数存于内存的真实地址。指令中地址码所表示的地址称为 (3) , 将此地址经过变换或运算而得到的操作数的真实地址称为 (4) , 相对于某一寄存器内容而言的距物理地址的差距值称为 (5) 。(5) 处填 ()。

A. 物理地址

B. 形式地址

C. 偏移地址

D. 间接地址

5. 设有 4 道作业, 它们的提交时间及执行时间如表所示。在单道程序环境下, 若采用先来先服务调度算法, 其平均周转时间为 (15) , 平均带权周转时间为 (16) 。(15) 处填 ()。

A. 0.95

B. 2.8

C. 2.45

D. 5.25

6. The notion of NP-completeness has provided a (66) mathematical definition for (67) intractability of NP problems. But this measure applies only to worst-case complexity. Being NP-complete does not (68) that a problem is intractable on the average case. Indeed, some NP-complete problems are (69) on average, though some may not be. Levin initiated the study of average-case intractability. He showed that a bounded tiling problem under a simple distribution is average-case NP-complete. Since then, several additional average-case NP-complete problems have been shown within Levin's (70) . This paper is intended to provide a comprehensive survey of average-case NP-complete problems that have been published so far, and the techniques of obtaining these results. (68) 处填 ()。

A. direct

B. navigate

C. indicate

D. show

7. 在书店受订管理中涉及到以下 3 个关系模式：书籍 Books (Bid, Bname, Price, Author, Publisher) 订单 Orders (Ordend, Orderdate, Cid) 订单明细 Orderlist (Orderid, Bid, Qty) 其中各属性的含义是：Bid 书籍编号，Price 单价，Author 作者，Publisher 出版商，Ordend 订单编号，Orderdate 下订日期，Cid 客户编号，Qty 数量。每张订单具有唯一的订单编号；每张订单编号中可包含多种书籍，但每种书籍的编号仅允许出现一次。则订单实体的主键是 (33)，订单明细实体的主键是 (34)。请将正面的 SQL 语句空缺部分补充完整。CREATE TABLE ORDERLIST (ORDERID CHAR (20), BD CHAR(6), QTY NUMERIC(9), (35) (ORDERID, BID), (36) (ORDERID) (37) (BID)) (35) 处填 ()。

A. PRIMARY KEY

B. FOREIGN KEY

C. FOREIGN KEY (ORDERID) REFERENCES ORDERS

D. FOREIGN KEY (BID) REFERENCES BOOKS

8. 在书店受订管理中涉及到以下 3 个关系模式：书籍 Books (Bid, Bname, Price, Author, Publisher) 订单 Orders (Ordend,

Orderdate, Cid) 订单明细 Orderlist (Orderid, Bid, Qty) 其中各属性的含义是: Bid 书籍编号, Price 单价, Author 作者, Publisher 出版商, Ordend 订单编号, Orderdate 下订日期, Cid 客户编号, Qty 数量。每张订单具有唯一的订单编号; 每张订单编号中可包含多种书籍, 但每种书籍的编号仅允许出现一次。则订单实体的主键是 (33), 订单明细实体的主键是 (34)。请将正面的 SQL 语句空缺部分补充完整。CREATE TABLE ORDERLIST (ORDERID CHAR (20), BD CHAR(6), QTY NUMERIC(9), (35) (ORDERID, BID), (36) (ORDERID) (37) (BID)) (36) 处填 ()。

- A. PRIMARY KEY
- B. FOREIGN KEY
- C. FOREIGN KEY (ORDERID) REFERENCES ORDERS
- D. FOREIGN KEY (BID) REFERENCES BOOKS

9. 设有 4 道作业, 它们的提交时间及执行时间如表所示。在单道程序环境下, 若采用先来先服务调度算法, 其平均周转时间为 (15), 平均带权周转时间为 (16)。(16) 处填 ()。

- A. 0.95
- B. 2.8
- C. 3.85

D. 5.25

10. The notion of NP-completeness has provided a (66) mathematical definition for (67) intractability of NP problems. But this measure applies only to worst-case complexity. Being NP-complete does not (68) that a problem is intractable on the average case. Indeed, some NP-complete problems are (69) on average, though some may not be. Levin initiated the study of average-case intractability. He showed that a bounded tiling problem under a simple distribution is average-case NP-complete. Since then, several additional average-case NP-complete problems have been shown within Levin's (70) . This paper is intended to provide a comprehensive survey of average-case NP-complete problems that have been published so far, and the techniques of obtaining these results. (69) 处填 ()。

A. difficult

B. easy

C. hard

D. intractable

11. 可行性分析的目的是在尽可能短的时间内用尽可能小的代价来确定问题是否有解。不属于在可行性分析阶段进行的工作是 (20) 。可行性分析不包括对待开发软件进行 (21) 分析。(20) 处填 ()。

- A. 研究目前正在使用的系统
- B. 根据待开发系统的要求导出新系统的逻辑模型
- C. 提供几个可供选择的方案
- D. 编制项目开发计划

12. 在书店受订管理中涉及到以下 3 个关系模式：书籍 Books (Bid, Bname, Price, Author, Publisher) 订单 Orders (Ordend, Orderdate, Cid) 订单明细 Orderlist (Orderid, Bid, Qty) 其中各属性的含义是：Bid 书籍编号，Price 单价，Author 作者，Publisher 出版商，Ordend 订单编号，Orderdate 下订日期，Cid 客户编号，Qty 数量。每张订单具有唯一的订单编号；每张订单编号中可包含多种书籍，但每种书籍的编号仅允许出现一次。则订单实体的主键是 (33) ，订单明细实体的主键是 (34) 。请将正面的 SQL 语句空缺部分补充完整。CREATE TABLE ORDERLIST (ORDERID CHAR (20) , BD CHAR(6) , QTY NUMERIC(9) , (35) (ORDERID, BID), (36) (ORDERID) (37) (BID)) (37) 处填 ()。

A. PRIMARY KEY

B. FOREIGN KEY

C. FOREIGN KEY (ORDERID) REFERENCES ORDERS

D. FOREIGN KEY (BID) REFERENCES BOOKS

13. The Turing machine is an abstract (71) of computer execution and storage introduced in 1936 by Alan Turing to give a mathematically precise definition of (72) . or 'mechanical procedure'. As such it is still widely used in theoretical computer science, especially in (73) theory and the theory of computation. The thesis that states that Turing machines indeed capture the informal notion of effective or mechanical method in logic and mathematics is known as Turing's thesis. Every Turing machine computes a certain (74) partial function over the strings over its alphabet. In that sense it behaves like a computer with a fixed program. However, as Alan Turing already described, we can encode the action table of every Turing machine in a string. Thus we might try to construct a Turing machine that expects on its tape a string describing an action table followed by a string describing the input tape, and then computes the tape that the encoded Turing

machine would have computed. As Turing showed, such a Turing machine is indeed possible and since it is able to simulate any other Turing machine it is called a (75) Turing machine. A universal Turing machine is Turing complete. It can calculate any recursive function, decide any recursive language, and accept any recursively enumerable language. According to the Church-Turing thesis, the problems solvable by a universal Turing machine are exactly those problems solvable by an algorithm or an effective method of computation, for any reasonable definition of those terms.

(72) 处填 ()。

- A. operation
- B. calculating
- C. algorithm
- D. mechanics

14. The notion of NP-completeness has provided a (66) mathematical definition for (67) intractability of NP problems. But this measure applies only to worst-case complexity. Being NP-complete does not (68) that a problem is intractable on the average case. Indeed, some

NP-complete problems are (69) on average, though some may not be. Levin initiated the study of average-case intractability, He showed that a bounded tiling problem under a simple distribution is average-case NP-complete. Since then, several additional average-case NP-complete problems have been shown within Levin's (70) . This paper is intended to provide a comprehensive survey of average-case NP-complete problems that have been published so far, and the techniques of obtaining these results. (70) 处填 ()。

- A. framework
- B. problem
- C. basis
- D. foundation

15. 可行性分析的目的在于在尽可能短的时间内用尽可能小的代价来确定问题是否有解。不属于在可行性分析阶段进行的工作是 (20) 。可行性分析不包括对待开发软件进行 (21) 分析。(21) 处填 ()。

- A. 技术可行性

B. 经济可行性

C. 操作可行性

D. 组织可行性

16. The Turing machine is an abstract (71) of computer execution and storage introduced in 1936 by Alan Turing to give a mathematically precise definition of (72) . or 'mechanical procedure'. As such it is still widely used in theoretical computer science, especially in (73) theory and the theory of computation. The thesis that states that Turing machines indeed capture the informal notion of effective or mechanical method in logic and mathematics is known as Turing's thesis. Every Turing machine computes a certain (74) partial function over the strings over its alphabet. In that sense it behaves like a computer with a fixed program. However, as Alan Turing already described, we can encode the action table of every Turing machine in a string. Thus we might try to construct a Turing machine that expects on its tape a string describing an action table followed by a string describing the input tape, and then computes the tape that the encoded Turing machine would have computed. As Turing showed, such a Turing

machine is indeed possible and since it is able to simulate any other Turing machine it is called a (75) Turing machine. A universal Turing machine is Turing complete. It can calculate any recursive function, decide any recursive language, and accept any recursively enumerable language. According to the Church-Turing thesis, the problems solvable by a universal Turing machine are exactly those problems solvable by an algorithm or an effective method of computation, for any reasonable definition of those terms.

(73) 处填 ()。

- A. intricacy
- B. complexity
- C. complicity
- D. difficulty

17. The Turing machine is an abstract (71) of computer execution and storage introduced in 1936 by Alan Turing to give a mathematically precise definition of (72) . or 'mechanical procedure'. As such it is still widely used in theoretical computer science, especially in (73) theory and the theory of computation. The thesis that

states that Turing machines indeed capture the informal notion of effective or mechanical method in logic and mathematics is known as Turing's thesis. Every Turing machine computes a certain (74) partial function over the strings over its alphabet. In that sense it behaves like a computer with a fixed program. However, as Alan Turing already described, we can encode the action table of every Turing machine in a string. Thus we might try to construct a Turing machine that expects on its tape a string describing an action table followed by a string describing the input tape, and then computes the tape that the encoded Turing machine would have computed. As Turing showed, such a Turing machine is indeed possible and since it is able to simulate any other Turing machine it is called a (75) Turing machine. A universal Turing machine is Turing complete. It can calculate any recursive function, decide any recursive language, and accept any recursively enumerable language. According to the Church-Turing thesis, the problems solvable by a universal Turing machine are exactly those problems solvable by an algorithm or an effective method of computation, for any reasonable definition of those terms.

(74) 处填 ()。

- A. fixed
- B. steady
- C. variational
- D. changeable

18. The Turing machine is an abstract (71) of computer execution and storage introduced in 1936 by Alan Turing to give a mathematically precise definition of (72) . or 'mechanical procedure'. As such it is still widely used in theoretical computer science, especially in (73) theory and the theory of computation. The thesis that states that Turing machines indeed capture the informal notion of effective or mechanical method in logic and mathematics is known as Turing's thesis. Every Turing machine computes a certain (74) partial function over the strings over its alphabet. In that sense it behaves like a computer with a fixed program. However, as Alan Turing already described, we can encode the action table of every Turing machine in a string. Thus we might try to construct a Turing machine that expects on its tape a string describing

an action table followed by a string describing the input tape, and then computes the tape that the encoded Turing machine would have computed. As Turing showed, such a Turing machine is indeed possible and since it is able to simulate any other Turing machine it is called a (75) Turing machine. A universal Turing machine is Turing complete. It can calculate any recursive function, decide any recursive language, and accept any recursively enumerable language. According to the Church-Turing thesis, the problems solvable by a universal Turing machine are exactly those problems solvable by an algorithm or an effective method of computation, for any reasonable definition of those terms.

(75) 处填 ()。

- A. universal
- B. special
- C. complete
- D. changeable

19. 嵌入到 Web 页中的 Flash 动画的扩展名是 (45) , 而用来向 Flash 应用程序添加交互性的语言是 (46) 。

- A. *.fla

以上内容仅为本文档的试下载部分，为可阅读页数的一半内容。如要下载或阅读全文，请访问：<https://d.book118.com/198126133027006135>