

Functional-Dependency Theory

- ◆ **We now consider the formal theory that tells us which functional dependencies are implied logically by a given set of functional dependencies.**
- ◆ **We then develop algorithms to generate lossless decompositions into BCNF and 3NF and higher normal form.**

Closure of a Set of Functional Dependencies

- ◆ Given a set F set of functional dependencies, there are certain other functional dependencies that are logically implied by F .
 - For e.g.: If $A \rightarrow B$ and $B \rightarrow C$, then we can infer that $A \rightarrow C$
- ◆ The set of **all** functional dependencies logically implied by F is the **closure** of F , we denote the *closure* of F by F^+ .

Closure of a Set of Functional Dependencies

- ◆ We can find F^+ , the closure of F , by repeatedly applying

Armstrong's Axioms:

- if $\beta \subseteq \alpha$, then $\alpha \rightarrow \beta$ (reflexivity 自反律)
 - if $\alpha \rightarrow \beta$, then $\gamma \alpha \rightarrow \gamma \beta$ (augmentation 增广律)
 - if $\alpha \rightarrow \beta$, and $\beta \rightarrow \gamma$, then $\alpha \rightarrow \gamma$ (transitivity 传递律)
- ◆ These rules are
 - sound (do not generate any incorrect functional dependencies)
 - complete (generate all functional dependencies that hold).

◆ $R = (A, B, C, G, H, I)$

$F = \{ A \rightarrow B$

$A \rightarrow C$

$CG \rightarrow H$

$CG \rightarrow I$

$B \rightarrow H\}$

◆ some members of F^+

● $A \rightarrow H,$

by transitivity from $A \rightarrow B$ and $B \rightarrow H$

● $AG \rightarrow I$

by augmenting $A \rightarrow C$ with G , to get $AG \rightarrow CG$
and then transitivity with $CG \rightarrow I$

● $CG \rightarrow HI$

by augmenting $CG \rightarrow I$ to infer $CG \rightarrow CGI$,
and augmenting of $CG \rightarrow H$ to infer $CGI \rightarrow HI$,
and then transitivity

Procedure for Computing F^+

- ◆ To compute the closure of a set of functional dependencies F :

$F^+ = F$

repeat

 for each functional dependency f in F^+

 apply reflexivity and augmentation rules on f

 add the resulting functional dependencies to F^+

 for each pair of functional dependencies f_1 and f_2 in F^+

 if f_1 and f_2 can be combined using transitivity

 then add the resulting functional dependency to F^+

until F^+ does not change any further

NOTE: We shall see an alternative procedure for this task later

Closure of Functional Dependencies (Cont.)

◆ Additional rules:

- If $\alpha \rightarrow \beta$ holds and $\alpha \rightarrow \gamma$ holds, then $\alpha \rightarrow \beta \gamma$ holds
(union)
- If $\alpha \rightarrow \beta \gamma$ holds, then $\alpha \rightarrow \beta$ holds and $\alpha \rightarrow \gamma$ holds
(decomposition)
- If $\alpha \rightarrow \beta$ holds and $\beta \rightarrow \delta$ holds, then $\alpha \rightarrow \delta$ holds
(pseudotransitivity)

The above rules can be inferred from Armstrong's axioms.

◆ **Example:** 已知关系模式R中

$U = \{A, B, C, D, E, G\}$,

$F = \{AB \rightarrow C, C \rightarrow A, BC \rightarrow D, ACD \rightarrow B,$

$D \rightarrow EG, BE \rightarrow C, CG \rightarrow BD, CE \rightarrow AG\}$,

判断 $BD \rightarrow AC$ 是否为F逻辑蕴含

◆ 解: 由 $D \rightarrow EG$ 知 $D \rightarrow E$, $BD \rightarrow BE$... ①

又知 $BE \rightarrow C$, $C \rightarrow A$ 所以 $BE \rightarrow A$, $BE \rightarrow AC$...

②

由①、②知, $BD \rightarrow AC$, 所以 $BD \rightarrow AC$ 被F所蕴含

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