

Outline

String Matching Problem

Hash Table

Knuth-Morris-Pratt (KMP) Algorithm

Suffix Trie

Suffix Array

String Matching Problem

String Matching Problem

- ▶ Given a text T and a pattern P , find all occurrences of P within T
- ▶ Notations:
 - n and m : lengths of P and T
 - Σ : set of alphabets (of constant size)
 - P_i : i th letter of P (1-indexed)
 - a, b, c : single letters in Σ
 - x, y, z : strings

Example

- ▶ $T = \text{AGCATGCTGCAGTCATGCTTAGGCTA}$
- ▶ $P = \text{GCT}$
- ▶ P appears three times in T

- ▶ A naive method takes $O(mn)$ time
 - Initiate string comparison at every starting point
 - Each comparison takes $O(m)$ time

- ▶ We can do much better!

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Hash Function

- ▶ A function that takes a string and outputs a number
- ▶ A good hash function has few collisions
 - i.e., If $x \neq y$, $H(x) \neq H(y)$ with high probability
- ▶ An easy and powerful hash function is a polynomial mod some prime p
 - Consider each letter as a number (ASCII value is fine)
 - $H(x_1 \dots x_k) = x_1 a^{k-1} + x_2 a^{k-2} + \dots + x_{k-1} a + x_k \pmod{p}$
 - How do we find $H(x_2 \dots x_{k+1})$ from $H(x_1 \dots x_k)$?

Hash Table

- ▶ Main idea: preprocess T to speedup queries
 - Hash every substring of length k
 - k is a small constant
- ▶ For each query P , hash the first k letters of P to retrieve all the occurrences of it within T
- ▶ Don't forget to check collisions!

Hash Table

- ▶ Pros:
 - Easy to implement
 - Significant speedup in practice
- ▶ Cons:
 - Doesn't help the asymptotic efficiency
 - ▶ Can still take $\Theta(nm)$ time if hashing is terrible or data is difficult
 - A lot of memory consumption

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